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(出版者 / Publisher)

法政大学経済学部学会

(雑誌名 / Journal or Publication Title)

経済志林 / The Hosei University Economic Review

(巻 / Volume)

57

(号 / Number)

2

(開始ページ / Start Page)

95

(終了ページ / End Page)

106

(発行年 / Year)

1989-06-15

(URL)

<https://doi.org/10.15002/00008503>

# A Direct Proof of Aumann and Maschlers' Theorem on The Nucleolus of A Bankruptcy Game

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## Abstract

An alternative proof of Aumann and Maschlers' theorem on the nucleolus of a Talmudic bankruptcy game is given directly from the definition of the nucleolus.

## 1. Introduction

The purpose of this note is to give a direct proof of Aumann and Maschlers' interesting theorem on a bankruptcy problem based on the Talmud [1]. This theorem states that the CG-consistent solution to a bankruptcy problem, which is defined after a Talmudic principle called by them the *contested garment principle*, is precisely the nucleolus of a game associated with the bankruptcy problem.

Their short and elegant proof makes use of theorems of cooperative game theory, e. g., [2], [3], [4], some of them being not so familiar to non-specialists. The proof is completed by showing that the kernel of the associated game consists of a single point, thereby establishing the identity of it and the nucleolus via the theorem of Schmeidler [4].

In contrast, in the proof to be given below, we use only the definition of the nucleolus [4], which makes the proof rather lengthy, yet direct, quite elementary and more easily accessible for non-specialists.

We give only the definitions and results that are necessary for our proof. For motivations and discussions on them, refer to Aumann and Maschler [1]. The proof is performed in a straightforward manner by first representing formally the CG-consistent solution which Aumann and Maschler have described in their theorem A, and then showing directly that no other solution can satisfy the requirement of the nucleolus.

## 2. Definitions and the Theorem

A bankruptcy problem is a pair  $(E; d)$  where  $E$  is the estate of a bankrupt, and  $d = (d_1, \dots, d_n)$ ,  $0 \leq d_1 \leq \dots \leq d_n$ , is the debts to  $n$  creditors  $1, \dots, n$ , satisfying  $0 \leq E \leq d_1 + \dots + d_n \equiv D$ . A solution to  $(E; d)$  is an  $n$ -tuple  $x = (x_1, \dots, x_n)$  of real numbers with  $x_1 + \dots + x_n = E$ .

A solution  $x$  is called *CG-consistent* if for all  $i \neq j$ ,  $(x_i, x_j)$  satisfies

$$x_i = (X_{ij} - (X_{ij} - d_i)_+ - (X_{ij} - d_j)_+)/2 + (X_{ij} - d_j)_+$$

and

$$x_j = (X_{ij} - (X_{ij} - d_i)_+ - (X_{ij} - d_j)_+)/2 + (X_{ij} - d_i)_+,$$

where

$$\begin{aligned} X_{ij} &\equiv x_i + x_j, \\ t_+ &\equiv \max(t, 0). \end{aligned}$$

Aumann and Maschler [1] have shown that every bankruptcy problem has a unique CG-consistent solution.

In this note, a *game* is a function  $v$  that associates a nonnegative real number  $v(S)$  with each subset  $S$  of  $N = \{1, \dots, n\}$ .  $N$  is the set of *players*, and  $S$  is called a *coalition*. It is assumed that  $v(\emptyset) = 0$ . A *payoff vector* is a vector  $x = (x_1, \dots, x_n)$  with  $x_1 + \dots + x_n = v(N)$ , where  $x_i$  represents a payoff to player  $i$ . An imputation is a payoff vector  $x$  satisfying  $x_i \geq v(\{i\})$  for all  $i \in N$ .

A *bankruptcy game* associated with a bankruptcy problem  $(E; d)$  is a game  $v_{E;d}$  defined by

$$v_{E;d}(S) = (E - d(N - S))_+ \text{ for each } S \subset N,$$

where

$$z(R) \equiv \sum_{i \in R} z_i$$

for any  $R \subset N$  and any vector  $z = (z_1, \dots, z_n)$ .

The nucleolus of a game  $v$  is an imputation  $x$  obtained as follows [4]. For a given imputation  $y$ , let  $\theta(y)$  be a vector in  $R^{2^n}$ , the  $2^n$ -dimensional Euclidean space, the components of which are the numbers  $v(S) - y(S)$  for all subsets  $S$  arranged in the non-increasing order, i. e.,  $\theta_1(y) \geq \theta_2(y) \geq \dots \geq \theta_{2^n}(y)$ . Then, an imputation  $x$  is called the nucleolus of  $v$  if for any imputation  $y \neq x$ ,

$$\theta_{i_0}(x) < \theta_{i_0}(y)$$

where

$$i_0 \equiv \min \{h | \theta_h(x) \neq \theta_h(y)\}.$$

It is well known that every game  $v$  has a unique nucleolus [4].

The number  $v(S) - y(S)$  is called the excess of coalition  $S$  with respect to  $y$ . Thus, the nucleolus has the meaning that it minimizes the maximal excess among the coalitions.

A striking result about the nucleolus of the bankruptcy game is that it is precisely the CG-consistent solution of the bankruptcy problem. Namely,

*Theorem* (Aumann and Maschler [1]). The CG-consistent solution of a bankruptcy problem  $(E, d)$  is the nucleolus of the game  $v_{E;d}$ .

To prove this theorem from the definition of the nucleolus, we need an explicit representation of the CG-consistent solution. Let  $x$  be the CG-consistent solution. Then, following the construction in Theorem A of Aumann and Maschler [1],  $x$  can be given as follows:

*Case (i)* If  $0 \leq E \leq nd_1/2$ , then

$$x_i = E/n \text{ for all } i=1, \dots, n.$$

*Case (ii)* For  $k=0, 1, 2, \dots, n-2$ , if

$$(D - \sum_{j=k+1}^n (d_j - d_{k+1}))/2 \leq E \leq (D - \sum_{j=k+2}^n (d_j - d_{k+2}))/2,$$

then

$$\begin{aligned} x_i &= d_i/2 \text{ for } i=1, \dots, k+1 \\ x_i &= c_{k+1} \text{ for } i=k+2, \dots, n, \end{aligned}$$

where

$$c_{k+1} = d_{k+1}/2 + \{E - (D - \sum_{j=k+1}^n (d_j - d_{k+1}))/2\} / (n - k - 1).$$

In this case we have

$$x_i \leq d_i/2, \text{ for } i=k+2, \dots, n.$$

To see this, put  $E = (D - \sum_{j=k+2}^n (d_j - d_{k+2}))/2$  for each  $k$  to obtain  $c_{k+1} \leq d_{k+2}/2$ .

*Case (iii)* For  $k=n-2, n-3, \dots, 1, 0$ , if

$$(D + \sum_{j=k+2}^n (d_j - d_{k+2}))/2 \leq E \leq (D + \sum_{j=k+1}^n (d_j - d_{k+1}))/2,$$

then

$$\begin{aligned} x_i &= d_i/2 \text{ for } i=1, \dots, k+1 \\ x_i &= d_i - b_{k+1} \text{ for } i=k+2, \dots, n, \end{aligned}$$

where

$$b_{k+1} = d_{k+1}/2 + \{(D + \sum_{j=k+1}^n (d_j - d_{k+1}))/2 - E\} / (n - k - 1).$$

In this case we have

$$x_i \geq d_i/2, \text{ for } i=k+2, \dots, n.$$

To see this, put  $E = (D + \sum_{j=k+2}^n (d_j - d_{k+2}))/2$  for each  $k$  to obtain  $b_{k+1} \leq d_{k+2}/2$ .

*Case (iv)* If  $D - nd_1/2 \leq E \leq D$ , then

$$x_i = d_i - (D - E)/n \text{ for all } i=1, \dots, n.$$

It will be convenient to note that the four cases are arranged

in the increasing order of  $E$  from 0 to  $D$ .

### 3. Proofs

Initially, we state four easy lemmas. The bankruptcy game  $v_{E;d}$  will be denoted simply by  $v$ .

*Lemma 1.* If  $E \leq (D - (d_n - d_{n-1}))/2$ , then  
 $v(\{i\}) = 0$  for all  $i = 1, \dots, n$ .

*Proof.* Note that  $D - d_n \geq d_{n-1}$ . Then,  
 $E \leq (D - d_n)/2 + d_{n-1}/2 \leq D - d_n$

Hence, for all  $i$ ,

$$\begin{aligned} 0 \leq v(\{i\}) &= \max\{0, E - D + d_i\} \\ &\leq \max\{0, E - D + d_n\} = v(\{n\}) = 0. \end{aligned}$$

*Lemma 2.* If  $(D + (d_n - d_{n-1}))/2 \leq E$ , then  
 $v(N - \{i\}) = E - d_i$  for all  $i = 1, \dots, n$ .

*Proof.*

$$E \geq d_n/2 + (D - d_{n-1})/2 \geq d_n.$$

Hence,  $E \geq d_i$  for all  $i = 1, \dots, n$ , which implies

$$v(N - \{i\}) = \max\{0, E - d_i\} = E - d_i, \text{ for all } i = 1, \dots, n.$$

*Lemma 3.* If

$$(D - \sum_{j=k+1}^n (d_j - d_{k+1}))/2 \leq E,$$

then

$$v(N - \{i\}) = E - d_i \text{ for all } i = 1, \dots, k+1.$$

*Proof.*

$$\begin{aligned} E &\geq D/2 - \sum_{j=k+1}^n (d_j - d_{k+1})/2 \\ &= (\sum_{j=1}^n d_j + \sum_{j=k+1}^n d_{k+1})/2 \geq 2d_{k+1}/2, \end{aligned}$$

because  $k \leq n-2$ . Hence, for all  $i=1, \dots, k+1$ ,

$$E-d_i \geq E-d_{k+1} \geq 0,$$

which implies

$$v(N-\{i\}) = \max\{0, E-d_i\} = E-d_i \text{ for all } i=1, \dots, k+1.$$

*Lemma 4.* If

$$E \leq (D + \sum_{j=k+1}^n (d_j - d_{k+1})),$$

then

$$v(\{i\}) = 0 \text{ for all } i=1, \dots, k+1.$$

*Proof.*

$$\begin{aligned} E &\leq D/2 + (D - \sum_{j=1}^n d_j - \sum_{j=k+1}^n d_{k+1})/2 \\ &= D - (\sum_{j=1}^n d_j + \sum_{j=k+1}^n d_{k+1})/2 \\ &\leq D - 2d_{k+1}/2 = D - d_{k+1}, \end{aligned}$$

because  $k \leq n-2$ . Hence, for all  $i=1, \dots, k+1$ ,

$$E - D + d_i \leq E - D + d_{k+1}$$

which implies

$$v(\{i\}) = 0 \text{ for all } i=1, \dots, k+1.$$

We now prove the theorem. Cases (i) and (iv) are proved before cases (ii) and (iii). In all proofs,  $S$  will stand for a nonempty, proper subset of  $N$ . The values  $v(S) - x(S)$  for  $S = \emptyset$  or  $N$  are always 0, so that they can be ignored.

*Case (i)* If  $0 \leq E \leq nd_1/2$ , then

$$x_i = E/n \text{ for all } i=1, \dots, n.$$

*Proof.* We show that if  $S \subset N$ ,  $S \neq N$  then

$$v(S) - x(S) \leq v(\{i\}) - x_i = -E/n \text{ for all } i=1, \dots, n.$$

Note that

$$E \leq nd_1/2 = (D - \sum_{j=1}^n (d_j - d_1))/2 \leq (D - (d_n - d_{n-1}))/2.$$

It then follows from Lemma 1 that

$$v(\{i\})=0 \text{ for all } i=1, \dots, n.$$

Then, if  $v(S)=0$ , there is a  $j$  such that

$$v(S)-x(S)\leq v(\{j\})-x_j=v(\{i\})-x_i=-E/n$$

for all  $i=1, \dots, n$ . If  $v(S)>0$ , then noting that  $x_i\leq d_i$  for all  $i=1, \dots, n$ , we have for some  $j$ ,

$$\begin{aligned} v(S)-x(S) &\leq v(N-\{j\})-x(N-\{j\})=-d_j+x_j \\ &=-d_j+E/n\leq -E/n \\ &=-x_i \text{ for all } i=1, \dots, n. \end{aligned}$$

Thus, the assertion is true. This also implies that  $x$  is an imputation.

Now, let  $y$  be any payoff vector with  $y\neq x$ . Then, for some  $i$ , we must have  $y_i<x_i$ . Hence,

$$v(\{i\})-y_i>v(\{i\})-x_i=-E/n \text{ for this } i,$$

which implies that  $y$  is not the nucleolus.

*Case (iv)* If  $D-nd_1/2\leq E\leq D$ , then

$$x_i=d_i-(D-E)/n \text{ for all } i=1, \dots, n.$$

*Proof.* We show that if  $S\subset N$ ,  $S\neq N$  then

$$v(S)-x(S)\leq v(N-\{i\})-x(N-\{i\})=-(D-E)/n$$

for all  $i=1, \dots, n$ . Note that

$$E\geq D-nd_1/2=(D+\sum_{j=1}^n(d_j-d_1))/2\geq(D+(d_n-d_{n-1}))/2.$$

It then follows from Lemma 2 that

$$v(N-\{i\})=E-d_i \text{ for all } i=1, \dots, n.$$

If  $v(S)>0$ , then noting that  $x_i\leq d_i$  for all  $i=1, \dots, n$ , we have for some  $j$ ,

$$\begin{aligned} v(S)-x(S) &\leq v(N-\{j\})-x(N-\{j\})=-d_j+x_j \\ &=-(D-E)/n=-d_i+x_i \\ &=v(N-\{i\})-x(N-\{i\}) \text{ for all } i=1, \dots, n. \end{aligned}$$

If  $v(S)=0$ , then for some  $j$  we have

$$v(S)-x(S)\leq v(\{j\})-x_j$$



$$\begin{aligned}
&= -d_j + (D - E)/n \leq -(D - E)/n \\
&= v(N - \{i\}) - x(N - \{i\}) \text{ for all } i=1, \dots, n.
\end{aligned}$$

Thus, the assertion is true and  $x$  is an imputation.

Now, let  $y$  be any payoff vector with  $y \neq x$ . Then, for some  $i$ , we must have  $y_i > x_i$ . Hence,

$$v(N - \{i\}) - y(N - \{i\}) > v(N - \{i\}) - x(N - \{i\}) \text{ for this } i,$$

which implies that  $y$  is not the nucleolus.

*Case (ii)* For  $k=0, 1, 2, \dots, n-2$ , if

$$(D - \sum_{j=k+1}^n (d_j - d_{k+1}))/2 \leq E \leq (D - \sum_{j=k+2}^n (d_j - d_{k+2}))/2,$$

then

$$\begin{aligned}
x_i &= d_i/2 \text{ for } i=1, \dots, k+1 \\
x_i &= c_{k+1} \text{ for } i=k+2, \dots, n,
\end{aligned}$$

where

$$c_{k+1} = d_{k+1}/2 + \{E - (D - \sum_{j=k+1}^n (d_j - d_{k+1}))/2\} / (n - k - 1)$$

*Proof.* Assume first that  $k \leq n-3$ . Then,  $E \leq (D - (d_n - d_{n-1}))/2$ . Hence, by Lemma 1,

$$v(\{i\}) = 0 \text{ for all } i=1, \dots, n.$$

Also, by Lemma 3,

$$v(N - \{i\}) = E - d_i \text{ for all } i=1, \dots, k+1.$$

We show that for each  $i=1, \dots, k$ , if  $S$  satisfies

$S \neq \{1\}$ ,  $S \neq N - \{1\}$ ,  $S \neq \{2\}$ ,  $S \neq N - \{2\}$ , ...,  $S \neq \{i\}$ , and  $S \neq N - \{i\}$ , then

$$\begin{aligned}
v(S) - x(S) &\leq v(N - \{i+1\}) - x(N - \{i+1\}) \\
&= v(\{i+1\}) - x_{i+1} \\
&= -d_{i+1}/2
\end{aligned} \tag{1}$$

Recall that  $x_i \leq d_i/2 < d_i$  for all  $i=1, 2, \dots, n$ . Then, for any such  $S$ , we have :

$v(S) > 0$  implies  $\exists h_{i+1} \neq 1, 2, \dots, i$  such that

$$\begin{aligned}
v(S) - x(S) &\leq v(N - \{h_{i+1}\}) - x(N - \{h_{i+1}\}) \\
&= -d_{h_{i+1}} + x_{h_{i+1}} \leq -x_{h_{i+1}},
\end{aligned}$$

and

$$\begin{aligned} v(S)=0 \text{ implies } \exists j_{i+1} \neq 1, 2, \dots, i \text{ such that} \\ v(S) - x(S) \leq v(\{j_{i+1}\}) - x_{j_{i+1}} \\ = -x_{j_{i+1}}. \end{aligned}$$

But, by Lemmas 3 and 1, we have

$$\begin{aligned} v(N - \{i+1\}) - x(N - \{i+1\}) &= -d_{i+1} + x_{i+1} \\ &= -d_{i+1}/2 - x_{i+1} \geq -x_{h_{i+1}} \end{aligned}$$

and

$$v(\{i+1\}) - x_{i+1} = 0 - x_{i+1} = -d_{i+1}/2 \geq -x_{j_{i+1}}.$$

Hence, (1) holds.

We next show that if  $S$  satisfies

$$S \neq \{1\}, S \neq N - \{1\}, \dots, S \neq \{k+1\}, S \neq N - \{k+1\},$$

then

$$\begin{aligned} v(S) - x(S) &\leq v(\{j\}) - x(\{j\}) \\ &= -c_{k+1} \leq -d_{k+1}/2 \text{ for all } j = k+2, \dots, n-1, n. \quad (2) \end{aligned}$$

This is because we have:

$$\begin{aligned} v(S)=0 \text{ implies } v(S) - x(S) &\leq v(\{j\}) - x_j \\ &= -x_j = -c_{k+1} \leq -d_{k+1}/2 \end{aligned}$$

and

$$\begin{aligned} v(S) > 0 \text{ implies } v(S) - x(S) &\leq v(N - \{j\}) - x(N - \{j\}) \\ &= -d_j + x_j \leq -x_j \leq -d_{k+1}/2. \end{aligned}$$

Combining (1) and (2), and noting that

$$v(\{1\}) - x_1 = -d_1/2 = v(N - \{1\}) - x(N - \{1\}),$$

we conclude that the first  $n$  greatest values of  $v(S) - x(S)$  can be arranged in the non-increasing order as

$$-d_1/2 \geq -d_2/2 \geq \dots \geq -d_{k+1}/2 \geq -c_{k+1} = \dots = -c_{k+1} \quad (3)$$

which also implies that  $x$  is an imputation.

Now, let  $y$  be any payoff vector with  $y \neq x$ , and let

$$i_0 = \min \{i | y_i \neq x_i\}.$$

Then, if  $i_0 \leq k+1$ , it follows from (1) that

$$v(N - \{i_0\}) - y(N - \{i_0\}) > -d_{i_0}/2$$

or

$$v(\{i_0\}) - y_{i_0} > -d_{i_0}/2.$$

Hence,  $y$  cannot be the nucleolus. If  $i_0 \geq k+2$ , then due to the assumption that  $k \leq n-3$ , there is another  $j_0 \geq k+2$  such that

$$y_{i_0} < x_{i_0} \text{ implies } y_{j_0} > x_{i_0}$$

and

$$y_{i_0} > x_{i_0} \text{ implies } y_{j_0} < x_{i_0}$$

Hence, we must have either

$$v(\{i_0\}) - y_{i_0} > -c_{k+1},$$

or

$$v(\{j_0\}) - y_{j_0} > -c_{k+1},$$

which implies that  $y$  is not the nucleolus.

When  $k = n-2$ , we have

$$x_i = d_i/2 \quad i=1, 2, \dots, n-1,$$

$$x_n = c_n \geq d_{n-1}/2$$

and (3) now becomes

$$-d_1/2 \geq -d_2/2 \geq \dots -d_{n-1}/2 \geq -c_n.$$

Note that  $i_0 < n$  by definition. Hence, it follows from (1) that either

$$v(\{i_0\}) - y_{i_0} > -d_{i_0}/2$$

or

$$v(N - \{i_0\}) - y(N - \{i_0\}) > -d_{i_0}/2,$$

which implies that  $y$  is not the nucleolus. This completes the proof.

*Case (iii)* For  $k = n-2, n-3, \dots, 1, 0$ , if

$$(D + \sum_{j=k+2}^n (d_j - d_{k+2}))/2 \leq E \leq (D + \sum_{j=k+1}^n (d_j - d_{k+1}))/2,$$

then

$$x_i = d_i/2 \text{ for } i=1, \dots, k+1$$

$$x_i = d_i - b_{k+1} \text{ for } i=k+2, \dots, n,$$

$$b_{k+1} = d_{k+1}/2 + \{D + \sum_{j=k+1}^n (d_j - d_{k+1}))/2 - E\} / (n - k - 1).$$

*Proof.* The proof is similar to case (ii). Assume first that  $k \leq n-3$ . Then,  $E \geq (D + (d_n - d_{n-1}))/2$ . Hence, by Lemma 2,

$$v(N - \{i\}) = E - d_i \text{ for all } i=1, \dots, n.$$

Also, by Lemma 4,

$$v(\{i\})=0 \text{ for all } i=1, \dots, k+1.$$

We show that for each  $i=1, \dots, k$ , if  $S$  satisfies

$S \neq \{1\}$ ,  $S \neq N - \{1\}$ ,  $S \neq \{2\}$ ,  $S \neq N - \{2\}$ , ...,  $S \neq \{i\}$ , and  $S \neq N - \{i\}$ , then

$$\begin{aligned} v(S) - x(S) &\leq v(N - \{i+1\}) - x(N - \{i+1\}) \\ &= v(\{i+1\}) - x_{i+1} \\ &= -d_{i+1}/2 \end{aligned} \quad (1')$$

Recall that  $d_i/2 \leq x_i \leq d_i$  for all  $i=1, 2, \dots, n$ . Then, for any such  $S$ , we have:

$$\begin{aligned} v(S) > 0 \text{ implies } \exists h_{i+1} \neq 1, 2, \dots, i \text{ such that} \\ v(S) - x(S) &\leq v(N - \{h_{i+1}\}) - x(N - \{h_{i+1}\}) \\ &= -d_{h_{i+1}} + x_{h_{i+1}} \\ &= -d_{h_{i+1}}/2 \text{ if } h_{i+1} \leq k+1 \\ &= -b_{k+1} \text{ if } h_{i+1} \geq k+2 \end{aligned}$$

and

$$\begin{aligned} v(S) = 0 \text{ implies } \exists j_{i+1} \neq 1, 2, \dots, i \text{ such that} \\ v(S) - x(S) &\leq v(\{j_{i+1}\}) - x_{j_{i+1}} \\ &= -x_{j_{i+1}} \end{aligned}$$

But, by lemmas 2 and 4, we have

$$\begin{aligned} v(N - \{i+1\}) - x(N - \{i+1\}) &= -d_{i+1} + x_{i+1} \\ &= -d_{i+1}/2 \geq -d_{i+2}/2 \geq \dots \\ &\geq -d_{k+1}/2 \geq -b_{k+1}, \end{aligned}$$

and

$$v(\{i+1\}) - x_{i+1} = 0 - x_{i+1} = -d_{i+1}/2 \geq -x_{j_{i+1}}.$$

Hence, (1') holds.

We next show that if  $S$  satisfies

$$S \neq \{1\}, S \neq N - \{1\}, \dots, S \neq \{k+1\}, S \neq N - \{k+1\},$$

then

$$\begin{aligned} v(S) - x(S) &\leq v(N - \{j\}) - x(N - \{j\}) \\ &= -b_{k+1} \leq -d_{k+1}/2 \text{ for all } j=k+2, \dots, n-1, n. \end{aligned} \quad (2')$$

This is because we have:

$$v(S) > 0 \text{ implies } v(S) - x(S) \leq v(N - \{j\}) - x(N - \{j\})$$

$$= -d_i + x_j = -b_{k+1} \leq -d_{k+1}/2.$$

and

$$\begin{aligned} v(S)=0 \text{ implies } v(S)-x(S) &\leq v(\{j\})-x_j \\ &= -x_j \leq -d_j + x_j \\ &= -b_{k+1} \leq -d_{k+1}/2 \end{aligned}$$

Combining (1') and (2'), and noting that

$$v(\{1\})-x_1 = -d_1/2 = v(N-\{1\})-x(N-\{1\}),$$

we conclude that the first  $n$  greatest values of  $v(S)-x(S)$  can be arranged in the non-increasing order as

$$-d_1/2 \geq -d_2/2 \geq \dots \geq -d_{k+1}/2 \geq -b_{k+1} = \dots = -b_{k+1} \quad (3')$$

which also implies that  $x$  is an imputation. The rest of the proof is almost the same to that of case (ii), so is omitted.

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